

Scanstud

Evaluating Static Analysis Tools

SYSTEMATIC THOUGHT LEADERSHIP FOR INNOVATIVE BUSINESS



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Mission statement

- Investigating the state of the art in static analysis

Project overview

- Practical evaluation of commercial static analysis tools for security
- Focus on C and Java
- Done in 2008
- Joint work with the Siemens CERT

1. Introduction
2. Test methodology
3. Test code
4. Experiences and lessons learned

- 1. Introduction**
2. Test methodology
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What we WON'T tell you:

- The actual outcome of the evaluation
- Even if we wanted, we were not allowed (NDAs and such)

But:

- We do not consider the precise results to be too interesting
 - An evaluation as ours only documents a snapshot
 - and is outdated almost immediately

However:

- We hopefully will give you a general feel what can be assessed in respect to the capabilities of static analysis tools

So, what will we tell you



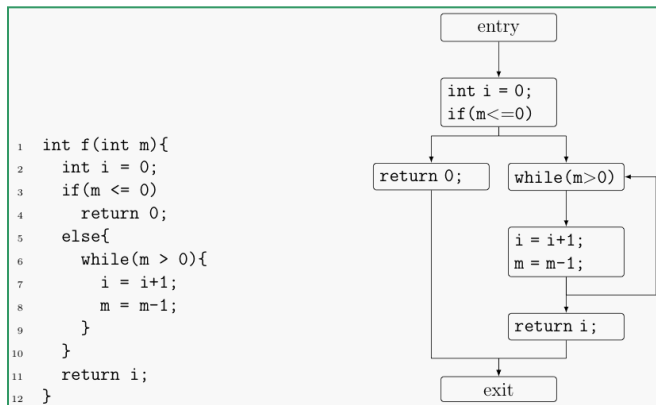
This talk is mainly about our evaluation methodology

- How we did it
- Why we did it this specific way
- General infos on the outcome
- Things we stumbled over

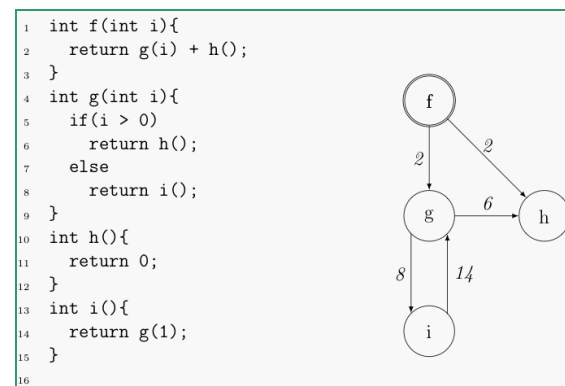
What makes a static analysis tool good?

It should find security problems

- Knowledge of different types of code based security problems
 - E.g., XSS, SQLi, Buffer Overflow, Format String problems...
- Language/Framework coverage
 - E.g., J2EE servlet semantics, <string.h>,...
- Understanding of flows
 - Control flow analysis (Loops invariants, integer ranges)
 - Data flow analysis (pathes from source to sink)



Control flow graph



Call graph

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- 2. Test methodology**
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Approaches

1. Use real world vulnerable software
2. Use existing or selfmade vulnerable application
 - Hacme, Web Goat, etc...
3. Create specific benchmarking suite

Our goal and how to reach it

- We want to learn a tool's specific capabilities
 - E.g., does it understand Arrays? Does it calculate loop invariants? Does it understand inheritance, scoping,...?
- Approaches 1. + 2. are not suitable
 - Potential side effects
 - more than one non-trivial operation in every execution path
- Writing custom testcode gives us the control that we need

However the other approaches are valuable too

Objectives

- Easy, reliable, correct, and iterative testcase creation
 - The actual test code should be
 - short
 - manual tested
 - as human readable as possible
- Defined scope of testcases
 - A single testcase should test only for one specific characteristic
- Automatic test-execution and -evaluation
 - Allows repeated testing and iterative testcase development
 - “neutral” evaluation

[Let's start at the bottom]

Approach

- Test-execution via batch-processing

Problem

- All tools behave differently

Solution

- Wrapper applications
 - Unified call interface
 - Unified XML-result format

Required

- Reliable mapping between alert and testcode

Approach

- One single vulnerability (or FP) per testcase
- Every testcase is hosted in an application of its own
- The rest of the application should otherwise be clean

Benefits

- Clear relation between alerts and testcases
 - Alert => the case was found / the FP triggered
 - No alert => the case was missed

Noise

- Even completely clean code can trigger warnings
 - The host-program may cause additional alerts
- How do we deterministically correlate scan-results to test-cases?
 - Line numbers are not always applicable.

Solution

- Result-Diff
 - Given two scan results it extracts the additional alerts
- Scan the host-program only (== the noise)
- Scan the host-program with injected testcase (== signal + noise)
- Diff the results (== signal)

Approach

- Separation between
 - **general support** code and
 - **test-specific** code (the actual vulnerabilities)

Benefit

- Support code is static for all testcases
- The actual testcase-code is reduced to the core of the tested property
 - Minimizes the code, reduces error-rate, increases confidentiality
 - Allows rapid testcase creation
 - Enables clear readability

Implementation

- Code generation
 - Host-program with defined insertion points
 - Testcode is inserted in the host-program

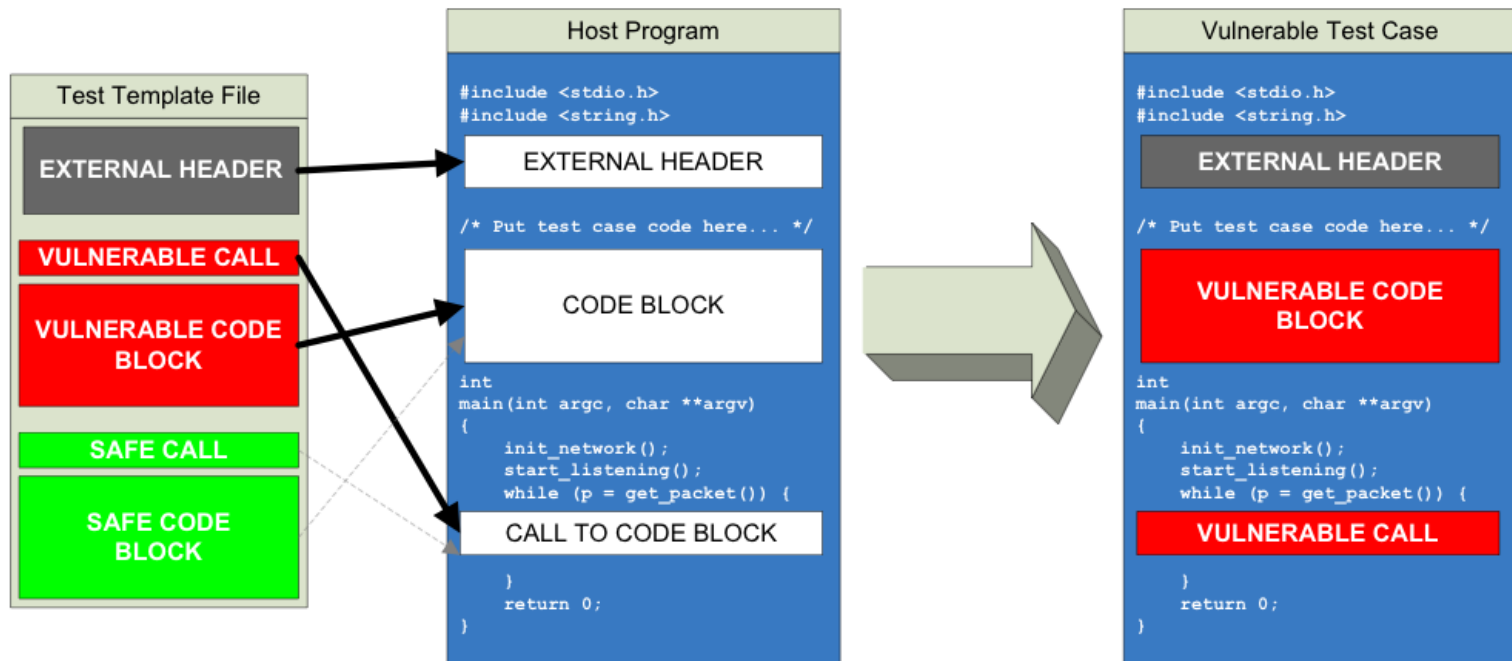
Insertion points in the host program

- Library includes, Global structures/data, function-call to the test function

The test-case is divided in several portions

- Each portion corresponds to one of the insertion points

A script merges the two files into one testcase



Example testcase(s): Buffer overflow



```
DESCRIPTION: Simple strcpy() overflow
ANNOTATION: Buffer Overflow [controlflow] []
```

EXTERNAL_HEADER:

```
#include <string.h>
```

```
VULNERABLE_CALL: %NAME(v)%(p);
```

VULNERABLE_EXTERNAL_CODE:

```
/* %DESCRIPTION(v)% */
void %NAME(v)%(char *p) {
    char buf[1024];
    strcpy(buf, p); /* %ANNOTATION(v)% */
}
```

```
SAFE_CALL: %NAME(s)%(p);
```

SAFE_EXTERNAL_CODE:

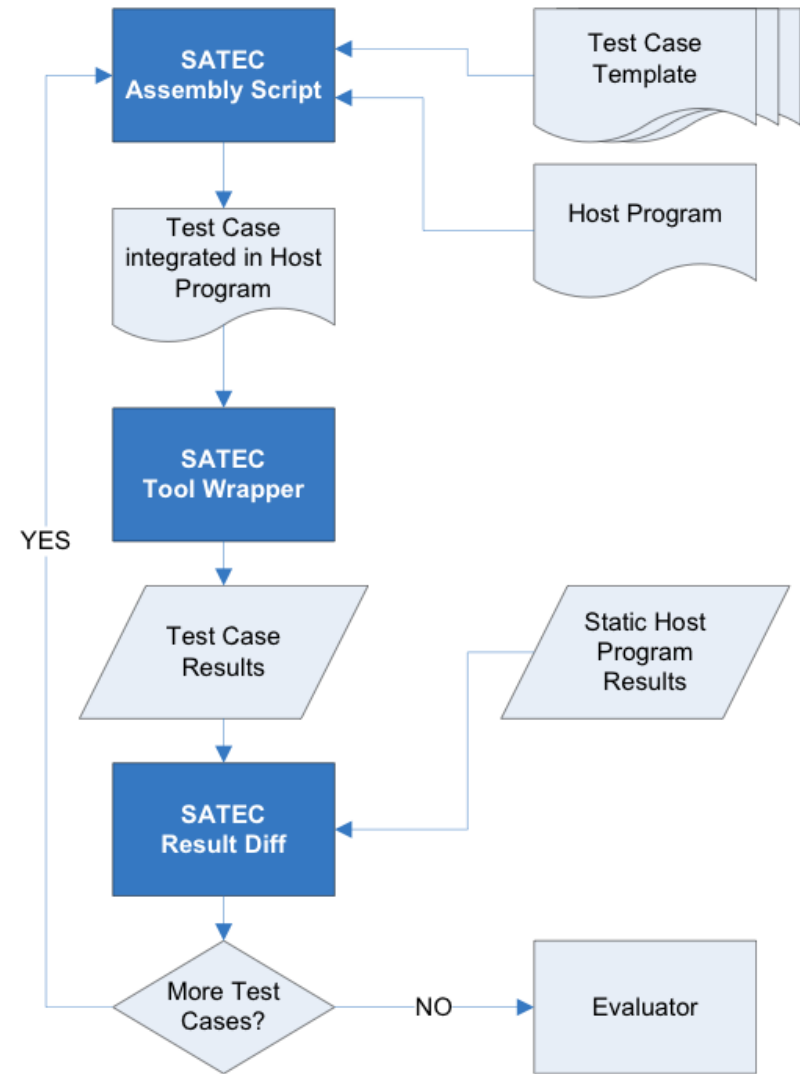
```
/* %DESCRIPTION(s)% */
void %NAME(s)%(char *p) {
    char buf[1024];
    if (strlen(p) >= sizeof(buf))
        return;
    strcpy(buf, p); /* %ANNOTATION(s)% */
}
```


Components

- Tool wrappers
- Host-program
- Test-cases
- Assembly script
- Result differ
- Evaluator

Putting it all together

- Creates test-code with the assembly-script
- Causes the wrapped tool to access the test-case
- Passes the test-result to result differ
- Dified-result and meta-data are finally provided to the Evaluator



Summary

- Applicable for all potential languages
- Applicable for all tools that provide a command-line interface
- Flexible
- Allows deterministic mapping code <--> findings

Fallback: Combined suite

- For cases where the tool cannot be wrapped
- All testcases are joined in one big application

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A **testcase** is the smallest unit in our approach

- Contains code which should probe for exactly one result
- Either “true vulnerability” or “false positive”

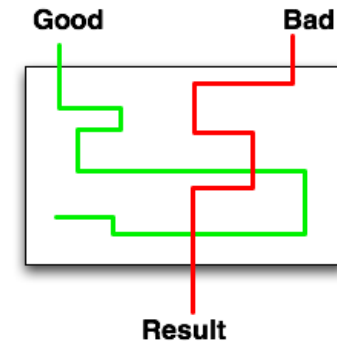
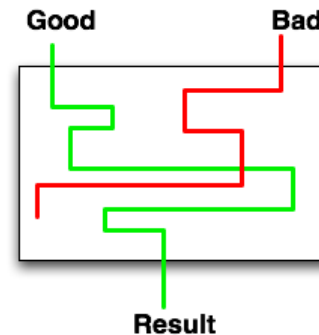
A **test** usually consists of two **testcases**

- a true vulnerability and
- a false positive
- Both testing the same characteristic

A test passed only if BOTH associated testcases have been identified correctly

Language features and control/data flow

- Two variables (“good”, “bad”) ⇒ The sources
 - Both are filled with user provided data
 - The “good” variable is properly sanitized
- One sink variable (“result”)
 - This variable is used to execute a security sensitive action
- Both variables are piped through a crafted control flow
- One of them is assigned to the result variable



Memory corruption

- Similar approach
- Instead of variables different sized memory regions are used

Host program

- All C test cases are hosted in a simple TCP server
- Listens on a port and waits for new clients
- Reads data from socket and passes pointer to test case
- Less than 100 LOC

The suite

- Emphasis on vulnerability types
- Around 116 single C test cases in total

Tests for, e.g.,

- Buffer overflows, unlimited/Off-by-one pointer loop overflows, integer overflows/underflows, signedness bugs, NULL pointer dereferences

Host program

- J2EE application with only one servlet
 - Provides: DB connection, framing HTML content, sanitizing,...

Vulnerability classes

- XSS, SQLi, Code Injection, Path Traversal, Response Splitting
 - ⇒ Emphasis on testing dataflow capabilities
- ~ 85 Java testcases in total
 - Ben Livshit's Stanford SecuriBench Micro was very helpful

Language features

- Library, inheritance, scoping, reflection, session storage

Tests

- Global buffers, array semantics, boolean logic, second order code injection, ...

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Market research: 12 potential candidates

■ Selection criteria:

- Maturity
- Is security a core-competence of the tool?
- Language support

⇒ Selection of 10 tools

⇒ After pre-tests 6 tools were chosen for further investigation

■ (no, we can't tell you which)

We have ~ 200 unique testcases

- How should the results be counted?

Observation

- If it aids the detection reliability, false positives are tolerable

Resulting quantification of the results

- Test passed: 3 Points
- False positive: 1 Point
- False negative: 0 Points

C Suite

Rank	Tool	Points
1.	Tool a.	72 / 168
2.	Tool b.	58 / 168
3.	Tool c.	56 / 168
4.	Tool d.	53 / 168
5.	Tool e.	50 / 168

Java Suite

Rank	Tool	Points
1.	Tool x.	89 / 147
2.	Tool y.	66 / 147
3.	Tool z.	58 / 147
4.	Tool v.	53 / 147

Categories covered by almost all tools:

- NULL pointer dereferences
- Double free's

Problem areas of most tools:

- Integer related bugs
 - Integer underflows / overflows leading to buffer overflows
 - Sign extension bugs
- Race conditions
 - Signals
 - `setjmp()` / `longjmp()`
- Non-implementation bugs
 - Authentication, Crypto, Privilege management, Truncation, ...

Strengths

- Within a function all tools possess good capabilities to track dataflows
- Besides that, the behaviour/capabilities are rather heterogeneous

Problem areas of most tools

- Global buffers
 - Especially if they are contained within a custom class
- Dataflow in and out of custom objects
 - E.g., our own linked list was too difficult for all tools

```
class Node {  
    public    String value;  
    public    Node   next;  
}
```

- Second order code injection

Buffer overflows 101:

- Most basic buffer overflow case?

```
strcpy()
```

- To our surprise, **3 out of 5** tools didn't report this!
 - Too obvious to report?
- One vendor was provided with this sample:

```
int main(int argc, char **argv) {  
    char buf[16];  
    strcpy(buf, argv[1])  
}
```

- Vendor response:

“argc/argv are not *modeled* to contain anything sensible.
We will eventually change that in the future.”

Buffer overflows 101:

- Another easy one:

```
gets (buf) ;
```

- Every tool must be finding that one!
 - Actually one tool didn't
- Vendor response:

“Ooops, this is a bug in our tool.”

More bugs:

- One tool didn't find anything in our "combined test case":

```
#include "testcase1.c"
#include "testcase2.c"
#include "testcase3.c"

int main(int argc, char **argv) {
    call_testcase1();
    call_testcase2();
    call_testcase3();
    return 0;
}
```

- Vendor response:

“#include'ed files are not analyzed *completely*.
Will be fixed in a future version.”

Let's sanitize some integers

- All tools allow the specification of sanitation functions
- So did Tool Y
- However the parameter for this function could only be
 - Int, float, ...
 - But not STRING!

Don't trust the servlet engine

- The J2EE host program writes some static HTML to the servlet response

```
PrintWriter writer = resp.getWriter();  
writer.println("<h3>ScanStud</h3>");
```

- Tool X warned "Validation needed"
 - (are you really sure you want your data there?)

One of the tools did not find a single XSS problem

- This surprised us, as the tool otherwise showed decent results
- Reason: We used the following code

```
PrintWriter writer = resp.getWriter();
```

- But the tool did not know “`getWriter()`”
- After replacing it with “`getOutputStream()`” XSS was found

Somewhat overeager

- Our SQLi tests exclusively used SELECT statements
- While detecting the vulnerability, the tool Z also warned
“stored XSS vulnerability”

A special price: The noisiest tool



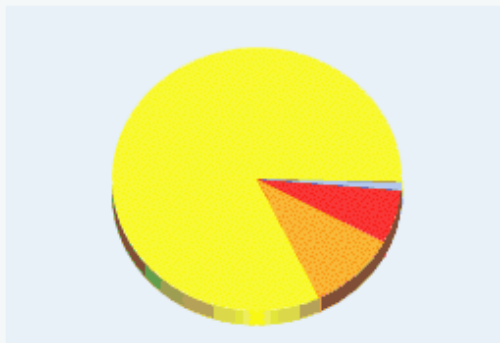
We had a tool in round one that did not understand neither C nor Java

- Therefore we started a C# benchmarking suite
- After three written testcases we did a first check
 - 2 XSS (vulnerable/safe), 1 SQLi (vulnerable)

484 Vulnerabilities

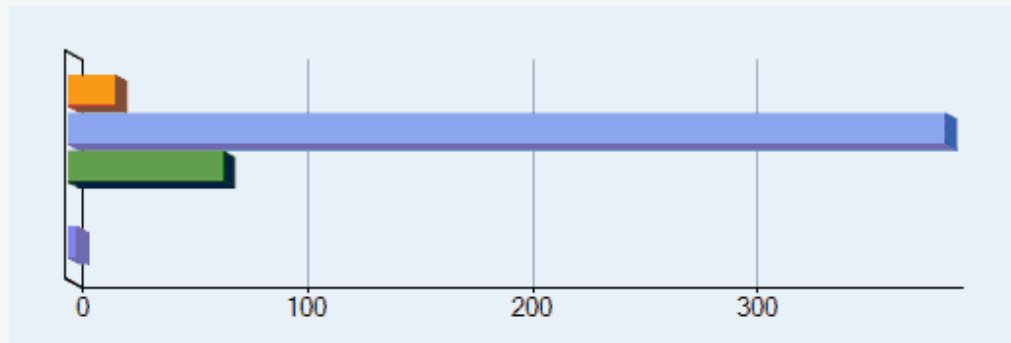
Analysis Completed

By Severity



● Critical	32
● Important	48
● Moderate	399
● Informational	5

By Category



■ Security Context	21
■ Insecure Coding Practices	390
■ Execution Errors	69
■ Application Integrity	0
■ Deployment Issues	4

Analysis duration: 00:01:48

Classes examined: 03

Page visits: 06

[View analysis transcript](#)

Questions?

□ Appendix

Pitfall

- Unbalanced creation/selection of testcases can introduce unsound results

Example

- Tool X is great but does not understand language feature Y
- Therefore all tests involving Y fail
- If there is an unbalanced amount of tests involving Y tool X has an unfair disadvantage

Solution: Categories and tags

- Categories: “controlflow”, “dataflow”, “language”,...
- Tags: All significant techniques within the testcase
 - Example: [cookies,conditional,loops]
- The it would be possible to see, that X always fails with Y

Vendor X:

- When there is a single path which includes an Array into a vulnerable data-flow, then the whole Array is tainted (even the safe values)
 - Underlying assumption: All elements of a linear data structure are on the same semantic level
 - This approach obviously breaks our test, to examine whether a tool understands Array semantics

Host program

- All C test cases are hosted in a simple TCP server
- Listens on a port and waits for new clients
- Accepts client connections
- Reads data from socket and passes pointer to test case
- Less than 100 LOC

Test cases

- Around 116 single C test cases in total
- 10 tests to determine the general *performance* of each tool
 - Arrays, loop constructs, structures, pointers, ...
- Rest of the test cases represent *real* vulnerabilities, which could be found in the wild

- Buffer overflows using simple unbounded string functions
 - strcpy, strcat, gets, fgets, sprintf, strvis, sscanf
- Buffer overflows using bounded string functions
 - snprintf, strncpy, strncat, memcpy
- Unlimited/Off-by-one pointer loop overflows
- Integer related bugs
 - Integer overflows / underflows
 - Sign extension
- Race conditions
 - Signals
 - setjmp()
 - TOCTTOU

- C operator misuse
 - sizeof(), assignment operator, octal numbers
- Format string issues
- NULL pointer derefs
- Memory management
 - Memory leaks
 - Double free's
- Privilege management
- Command injection
 - popen(), system()

The SATEC file format

- Each test is kept in a separate file
- The test is described using the following keywords
 - NAME (automatically generated from filename)
 - DESCRIPTION
 - ANNOTATION
- Two code blocks
 - VULNERABLE_EXTERNAL_CODE
 - SAFE_EXTERNAL_CODE
- Two calls, into the code blocks
 - VULNERABLE_CALL
 - SAFE_CALL
- Keyword expansion is possible

Example: T_001_C_XSS.java

DESCRIPTION: Very basic XSS

ANNOTATION: XSS [basic] []

VULNERABLE_CALL:

```
new %NAME(v)%().doTest(req, resp); // inserted by satec
```

SAFE_CALL:

```
new %NAME(s)%().doTest(req, resp); // inserted by satec
```

VULNERABLE_EXTERNAL_CODE:

```
class %NAME(v)% extends scanstudTestcase {
```

```
    public void doTest(HttpServletRequest req, HttpServletResponse resp){
```

```
        PrintWriter writer = resp.getWriter();
```

```
        String value = req.getParameter("testpar");
```

```
        writer.println("<h3>" + value + "</h3>"); // %ANNOTATION(v)%
```

```
    }
```

```
}
```

SAFE_EXTERNAL_CODE:

```
class %NAME(s)% extends scanstudTestcase {
```

```
    public void doTest(HttpServletRequest req, HttpServletResponse resp){
```

```
        PrintWriter writer = resp.getWriter();
```

```
        String value = HTML Encode(req.getParameter("testpar"));
```

```
        writer.println("<h3>" + value + "</h3>"); // %ANNOTATION(s)%
```

```
    }
```

```
}
```